

# Information-Theoretic Security

EE 25N, Science of Information

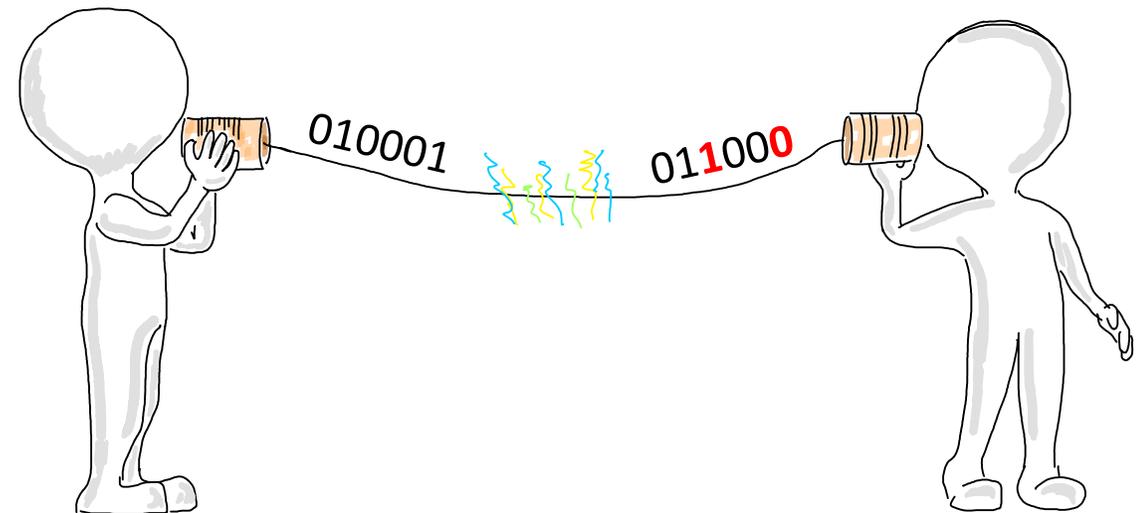
11/08/2018

Ziv Goldfeld

# Past Lectures - Communication Despite Noise

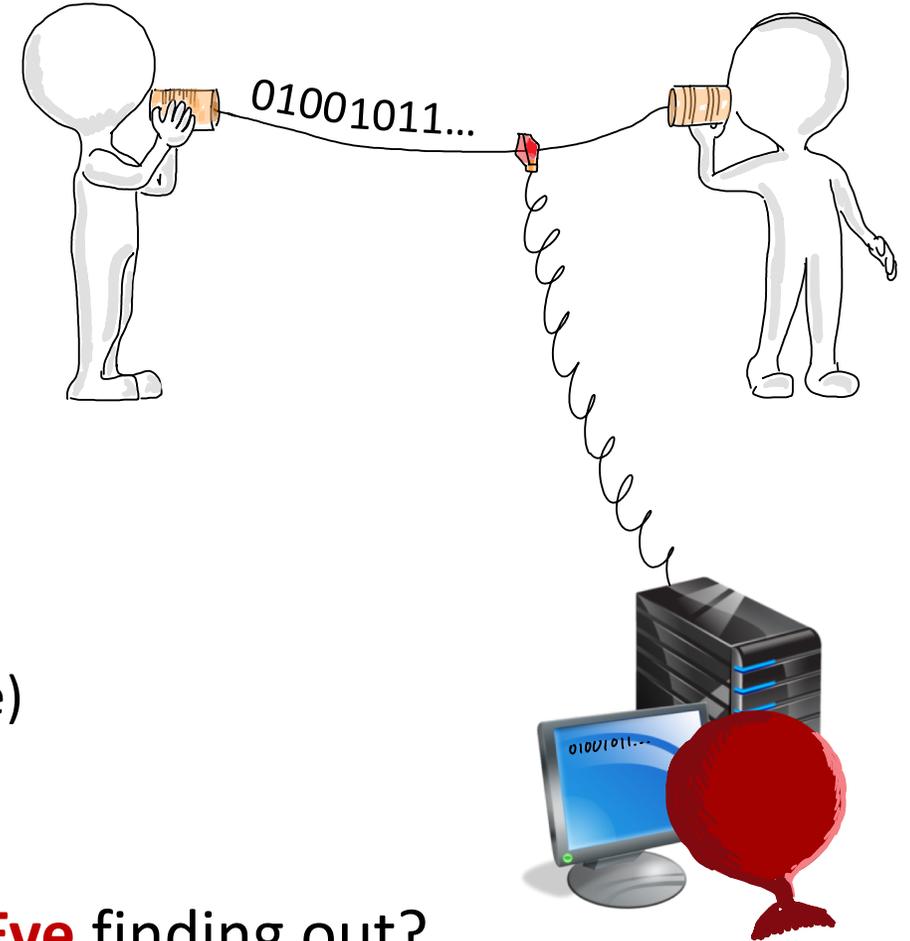
- Alice the sender and Bob the receiver
- Communicate a **message** as a strings of 0's and 1's (bits)
- Use **longer** bit strings (**codewords**) to protect **message** carrying bits
- Agree on a strategy beforehand:
  - 1) Set of codewords used
  - 2) Encoding for Alice
  - 3) Decoding for Bob

⇒ **Error correcting codes** for reliability



# Communication under Eavesdropping

- **Alice** and **bob** wish to communicate
- Channel is noiseless
- But **Eve** taps their line
- **They** don't want **Eve** to decipher their chat



- Assumptions on **Eve**:
  - 1) She sees their transmitted bit string
  - 2) She knows their communication strategy (aka code)
  - 3) She has an extremely powerful computer

**Q:** Can **Alice** send **Bob** a secret message without **Eve** finding out?

**A:** Not without an additional recourse!

Resource 1: Pre-Eve Secret

# Simple Case Study

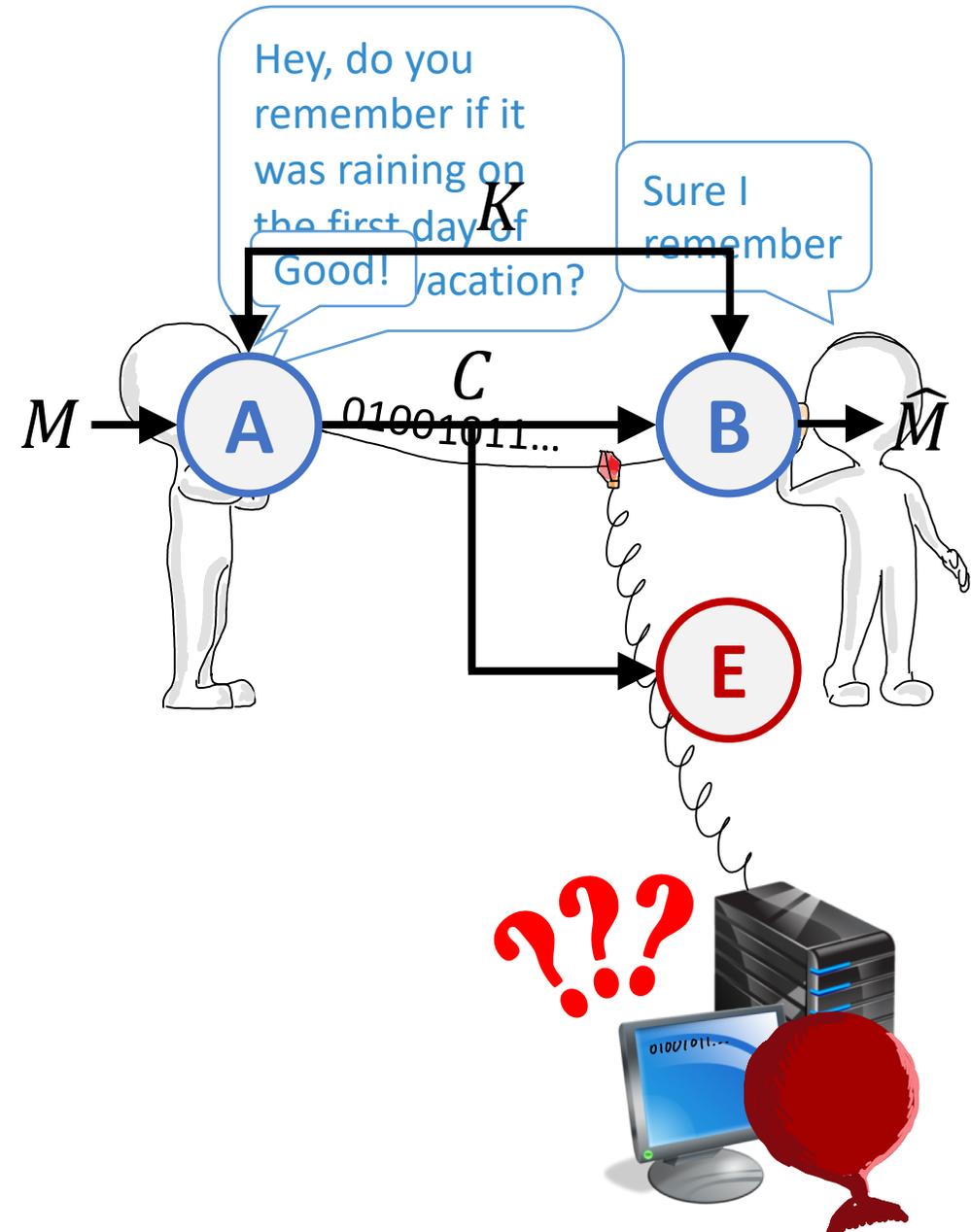
- **Alice** sends **Bob** a bit  $M \in \{0,1\}$   
Bit probability:  $P_M(0) = P_M(1) = \frac{1}{2}$
- They share a secret **Eve** has no access to  
 $\Rightarrow$  **Resource:** 1 secret bit  $K \in \{0,1\}$

Formally:

**Alice:**  $(M, K) \rightarrow C$

**Bob:**  $(C, K) \rightarrow \hat{M}$

**Eve:** Intercepts  $C$  and tries to figure out  $M$



# Simple Case Study – Modeling Eve

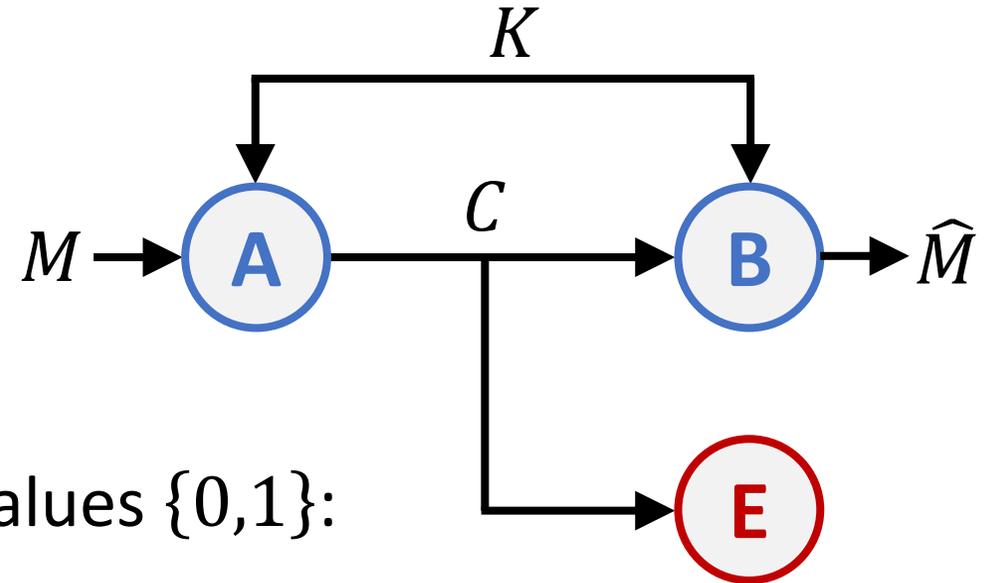
**Q1:** How to model **Eve's** perception of  $K$ ?

- Knows  $K$  is being used
- Doesn't know its value

⇒ **Eve** has a **guessing probability** over  $K$ 's values  $\{0,1\}$ :

- Doesn't have a clue:
- Knows something:

**Q2:** Which kind of secret should **Alice** and **Bob** favor?



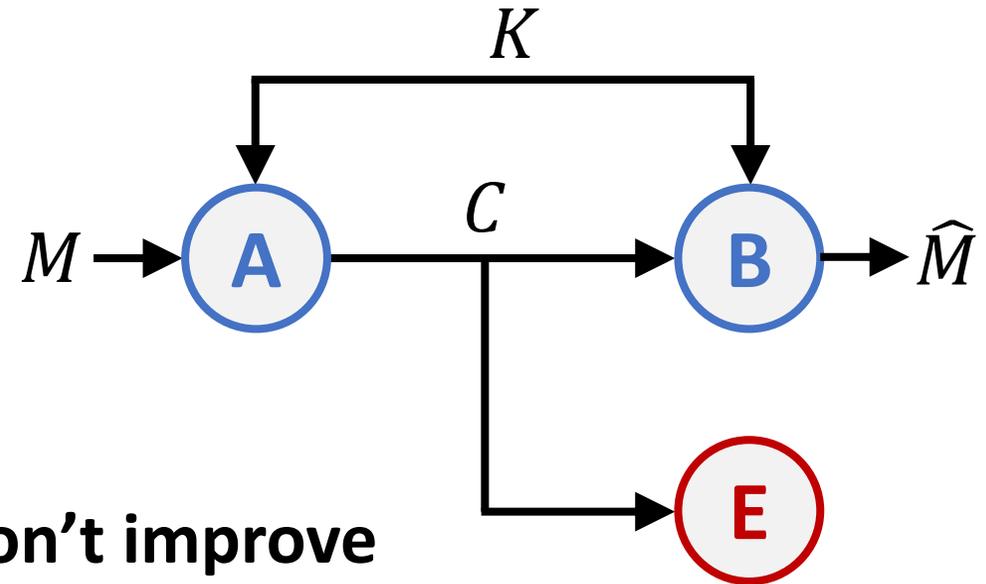
# Simple Case Study – Modeling Security

**Q3:** What does it mean to secure  $M$ ?

- Pre-transmission:  $P_M(0) = P_M(1) = \frac{1}{2}$

- **Eve** tries to recover  $M$  from  $C$

$\Rightarrow M$  is secure if **after** seeing  $C$  **Eve's** odds **don't improve**



**Goal:** Design functions for **Alice** and **Bob** such that:

- **Bob** can decode  $M$  from  $(C, K)$

- **Eve's** best guess of  $M$  **after** seeing  $C$  is still 50/50

# Simple Case Study – Binary Operations

- Assume  $M$  and  $K$  are both symmetric (50/50)
- **Alice** gets  $C$  via binary operation on  $(M, K)$
- Possible binary operations:

## OR

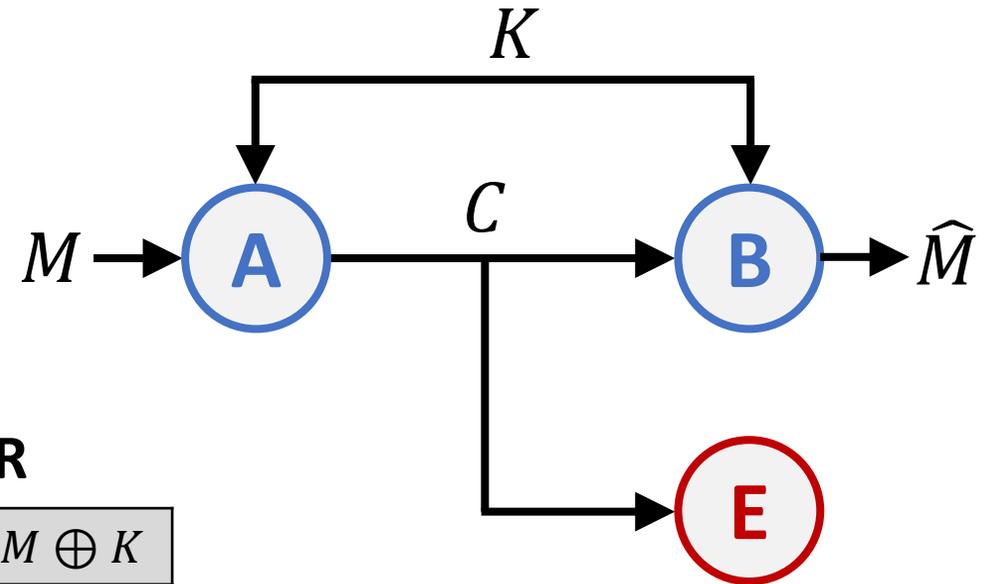
$M$	$K$	$M + K$
0	0	0
0	1	1
1	0	1
1	1	1

## AND

$M$	$K$	$M \cdot K$
0	0	0
0	1	0
1	0	0
1	1	1

## XOR

$M$	$K$	$M \oplus K$
0	0	0
0	1	1
1	0	1
1	1	0



**Q4:** Which binary operation is better for secrecy?

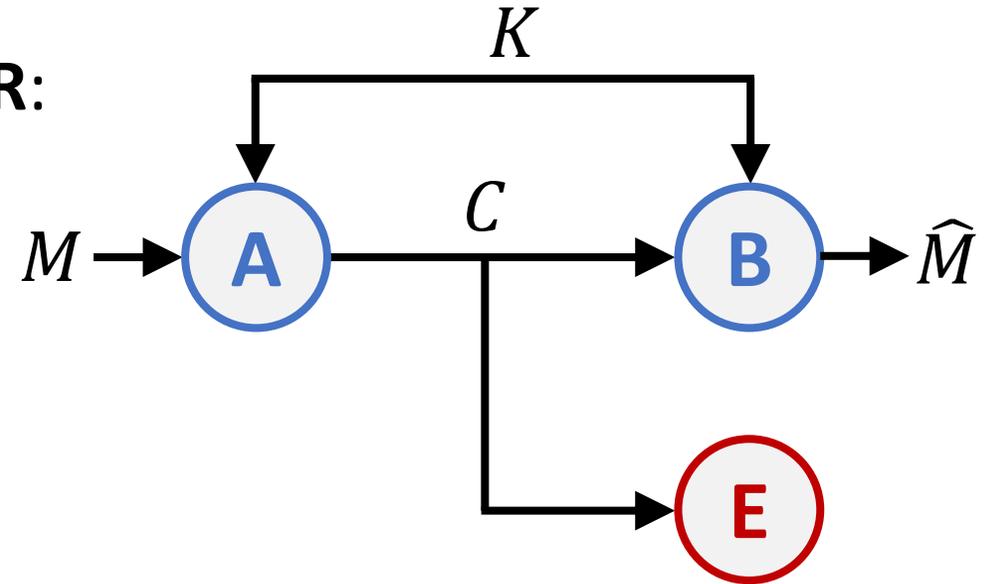
# Simple Case Study – Reliability & Optimality

⇒ Best function for symmetric  $(M, K)$  is **XOR**:

- **Eve**'s best guess after seeing  $C$  is 50/50

- **Same** odds like **before** seeing  $C$

⇒ Information-theoretic security ✓



**Q4:** Can **Bob** decode an **XOR**-based transmission? ✓

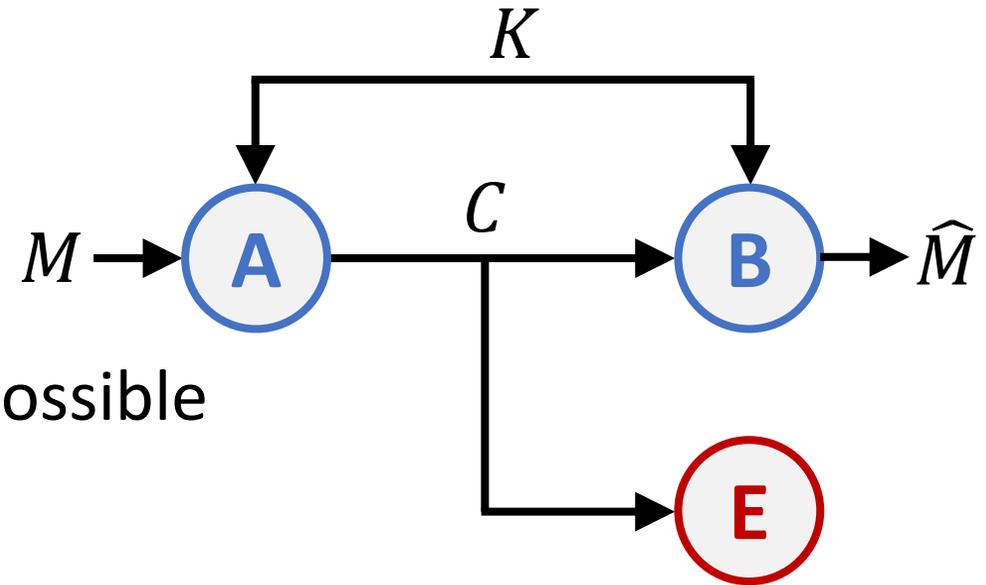
**Q5:** Can **OR** or **AND** operations be used for communication only? ✗

**Symmetry is Crucial:** Asymmetric keys can't achieve security with **XOR**

# Simple Case Study – General Claim

## One-Time Pad:

- $m$  messages bits and  $k$  key bits
- All bits are equiprobable
- $k^*$  = least  $k$  s.t. secure communication is possible



## Shannon (1949):

Achieving reliability & information-theoretic security over the OTP is:

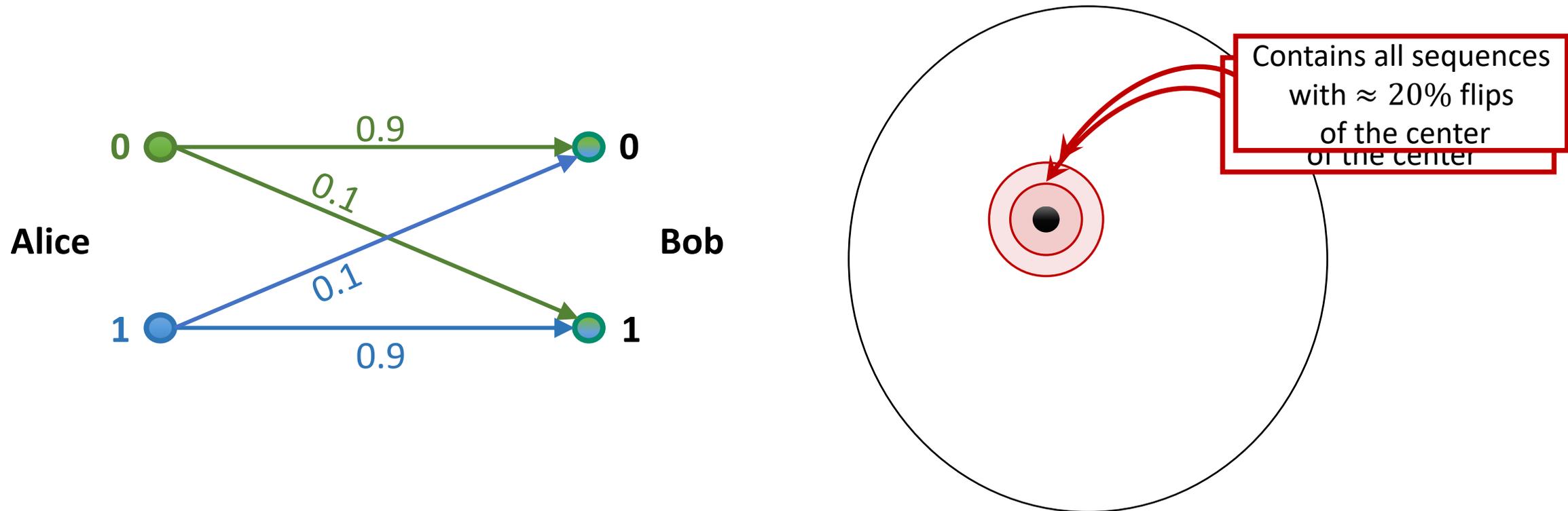
- 1) **possible** using exactly  $m$  key bits
- 2) **impossible** using less than  $m$  key bits

$$k^* = m$$

# Resource 2: Noise

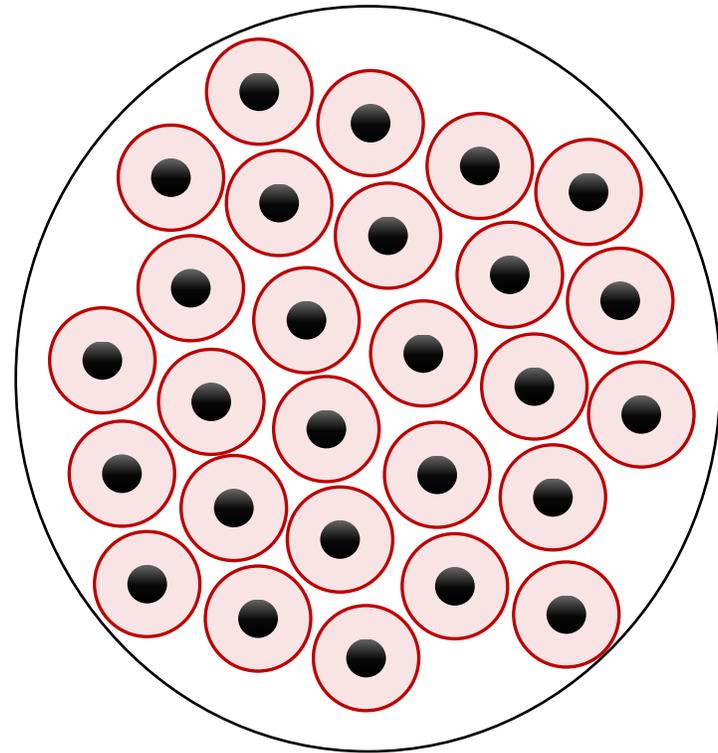
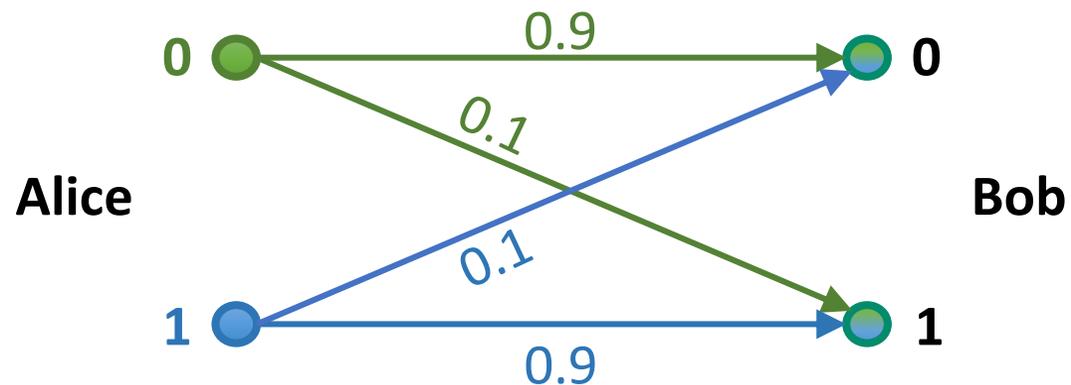
# Noisy Channel

- In most real-world systems we don't exactly know the number of bit flips
- Common mode of operation is to model noise probabilistically



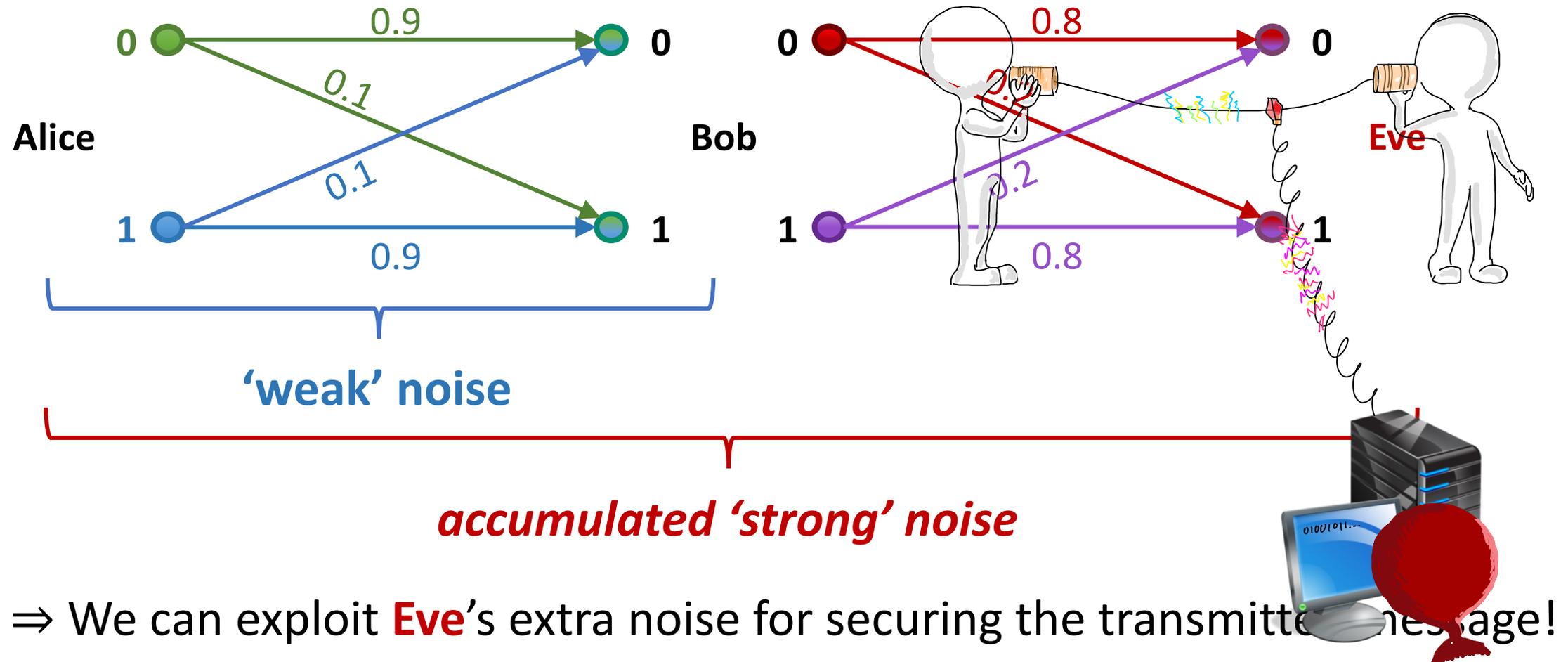
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# Wiretap Channel

Noisy communication channel with an eavesdropper



# Information-Theoretic Security Research

## Many interesting research questions:

- 1) Key agreement over noisy channels
- 2) Active Adversaries
  - **Eve** not only overhear the transmission but can influence the channel
  - Has a set of possible actions **Alice** and **Bob** know
  - They don't know which action is chosen  $\Rightarrow$  Ensure security versus all actions!
- 3) Covert Communication:
  - Communicate without **Eve** noticing
- 4) Many many many many more...